New Metrics for Reputation Management in P2P Networks

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Traditional Peer to Peer Networks

A type of network in which each workstation has equivalent capabilities and responsibilities. This differs from client/server architectures, in which some computers are dedicated to serving the others.

(Webopedia)™

P2P networks features

- ☑ Resource sharing: bandwidth, storage space, and computing power
- ☑ Information sharing
- ☑ Lack of central authority
- ☒ Lack of guarantee and certification of the shared resources

Downside

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The open and anonymous nature of P2P networks makes it difficult to calculate reliable quality metrics for peers and objects

Reputation management

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- Describe the performance of peers in the network
- Describe how reliable they are

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Such mechanisms should be robust against malicious peers.

Starting point

EigenTrust

We start with EigenTrust [Kamvar et al., 2003], an algorithm designed for reputation management in file sharing application over p2p networks. The main idea is to combine this algorithm with metrics of reputation computed using techniques recently introduced for detecting and demoting Web Spam.

Contribution

- We adapt Truncated PageRank [Becchetti et al., 2006], Estimation of Supporters [Palmer et al., 2002] and BadRank in reputation management
- We introduce a number of new threat models
- We test existing and new threat models in a simulated environment
- We show that our combined approaches perform better than EigenTrust alone in reducing the amount of inauthentic downloads

P2P networks and reputation
Preliminaries
Threat models
Metrics
Evaluation

EigenTrust

Applications of EigenTrust for reputation management

P2P networks (using a DHT to record transaction outcomes – never allow a peer to do its own evaluation), but also online communities

EigenTrust

Definition of local trust in EigenTrust

We define a local trust value s_{ij} as

$$s_{ij} = sat(i,j) - unsat(i,j).$$

In order to avoid malicious peers to assign arbitrarily high local trust values, it is necessary to normalize them. The normalized local trust value is c_{ij} is defined as follows:

$$cij = \frac{max(s_{ij},0)}{\sum_j max(s_{ij},0)}.$$

EigenTrust

Hypothesis

Peers who are honest about the files they provide are also likely to be honest in reporting their local trust values.

EigenTrust

Global trust

The idea of transitive trust, inspired by PageRank [Page et al., 1998], leads to a system where trust values propagate through paths along the network

PageRank

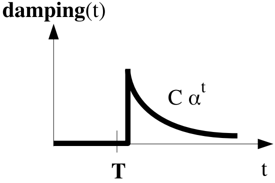
PageRank can be expressed as a weighted summation of paths of varying lengths

$$S = \sum_{t=0}^{\infty} \frac{\mathrm{damping}(t)}{N} P^t \ .$$

t: the lengths of the paths.
damping(t): decreasing function of t.
P: row-normalized citation matrix

Truncated PageRank

Proposed in [Becchetti et al., 2006]. Idea: reduce the direct contribution of the first levels of links:



$$\text{damping}(t) = \begin{cases} 0 & t \leq T \\ C\alpha^t & t > T \end{cases}$$

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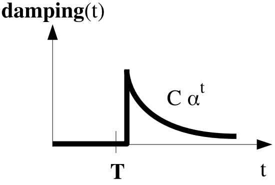

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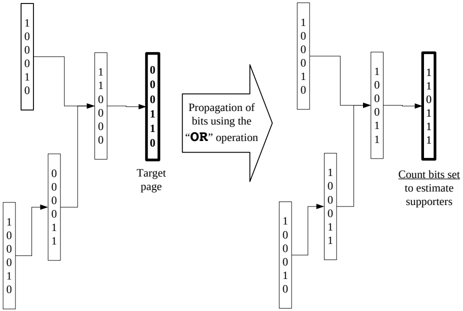
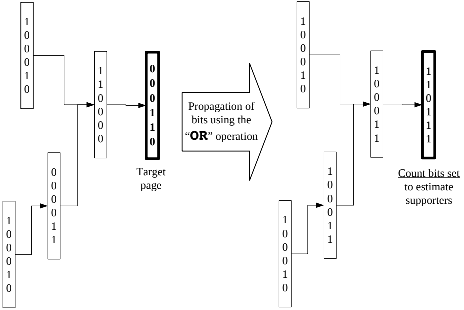
☑ No extra reading of the graph after PageRank

Estimation of supporters



Estimation of supporters



[Becchetti et al., 2006] shows an improvement of ANF algorithm [Palmer et al., 2002] based on probabilistic counting [Flajolet and Martin, 1985]. After d iterations, the bit vector associated to any page x provides information about the number of supporters of x at distance $\leq d$.
This algorithm can be used to estimate the number of different peers contributing to the ranking of a given peer.

BadRank

If a page links to another page with a high BadRank, then also this page should be considered a page with negative characteristics. The difference with respect to PageRank is that BadRank is not based on the evaluation of inbound links of a web page but on its outbound links.

$$br(i) = d \sum_{i \to j} \frac{br(j)}{indeg(j)} + (1-d)e(i)$$

computed on the graph of negative evaluations

Network Models

Transaction Network

A link from a node (peer) i to a node j is inserted every time i downloads a file from j. Each link is weighted with a positive value if the downloaded file was authentic, negative otherwise.

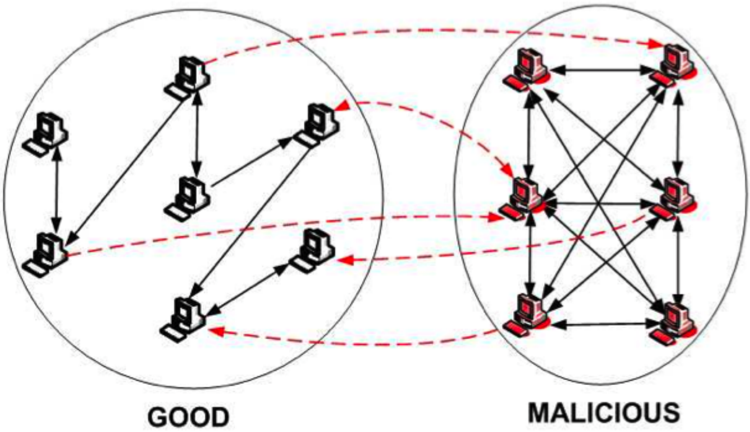
Positive Opinion Network

A link is inserted from a node i to a node j only after the download of authentic files.

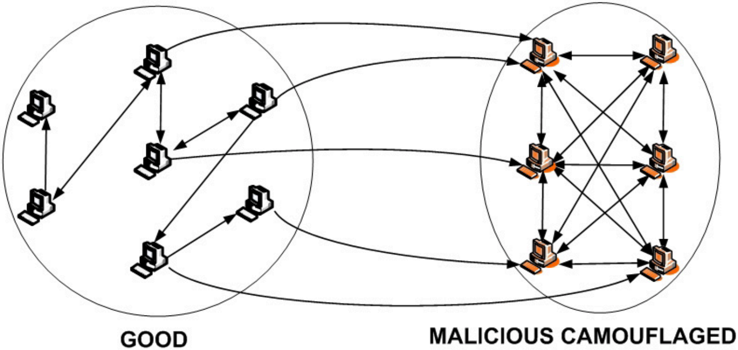
Inverse Network

The transpose of the positive opinion network.

Threat Model A (individuals) and B (collective)

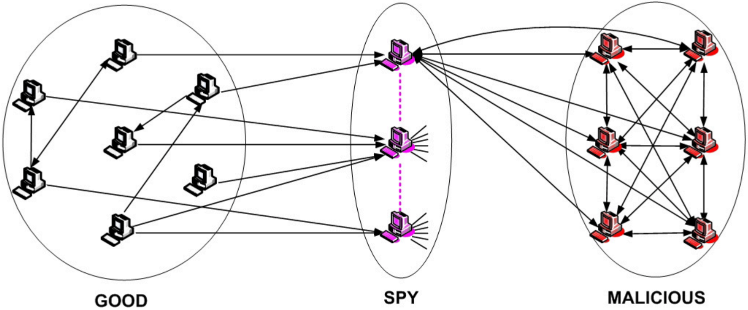


Threat Model C - collectives with camouflage



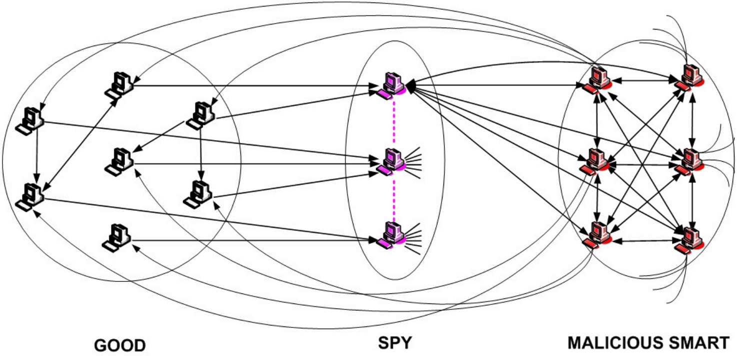
They provide good files sometimes

Threat Model D



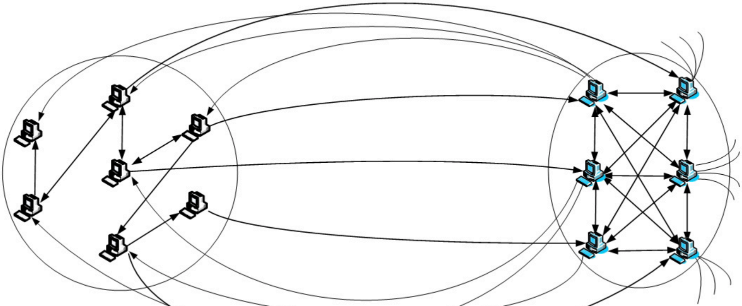
Have a set of nodes providing good ratings for them

Threat Model G - malicious smart model



Sometimes give ratings to the rest of the network

Threat Model H - malicious smart model with camouflage



Sometimes provide authentic files and ratings to the rest of the network

Eigentrust with Inverse Eigentrust - Model D

Encourage peers to provide ratings about other peers

Require: EigenTrust score vector ET, Inverse EigenTrust score vector I

if $I[i] > 0$ **then**
  return $ET[i]$
else
  return 0
end if

Eigentrust with Inverse Eigentrust - Model G

Encourage peers to provide many ratings about other peers

Require: EigenTrust score vector ET, Inverse EigenTrust score vector I, threshold $tr = \sum_i \frac{ET[i]}{N}$

if $I[i] \geq tr$ **then**
  return $ET[i]$
else
  return 0
end if

EigenTrust with Truncated PageRank

Malicious peers receive positive values from the other members of the coalition (malicious and spy). This means that the most of the *trust mass* is propagated starting from nodes at few hops of distance.

Require: Eigentrust score vector ET, Truncated PageRank vector P, threshold tr

if $P[i] \geq tr$ **then**
    return $ET[i]$
else
    return 0
end if

EigenTrust with Estimation of Supporters

Malicious peers supporters necessarily belong to the same coalition. This means that a malicious peer obtain an high reputation because of the great number of supporters at short distance from it.
The Bit Propagation algorithm can be used to perform an analysis of the connectivity of the transition network in order to detect local anomalies.

Require: EigenTrust score vector ET, Bit Propagation vector BP, threshold tr

```
if BP[i] ≥ tr then
    return ET[i]
else
    return 0
end if
```

Badness

Propagating badness

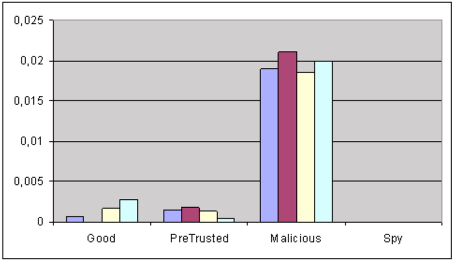
If i trusts j and j distrusts k then, with high probability, also i should regard k as untrustworthy. We can define the **Global Badness** as:

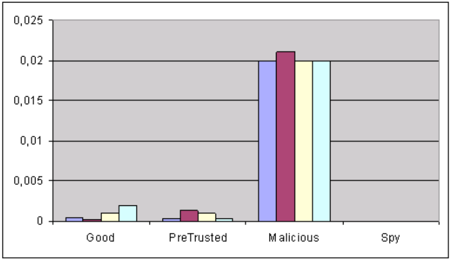
$$\mathbf{negT} = \mathrm{D}^{\top}\mathbf{T}$$

where D is the normalized negative opinion matrix and $\mathbf{T}$ is the EigenTrust Rank. Each peer i has a global Badness given by

$$\mathbf{negT_i} = \sum_{j=1}^{n} negC_{ji} \times \mathbf{T_j}$$

Average BadRank for models A-D





Average BadRank after 25 and 50 cycles.

Dishonesty

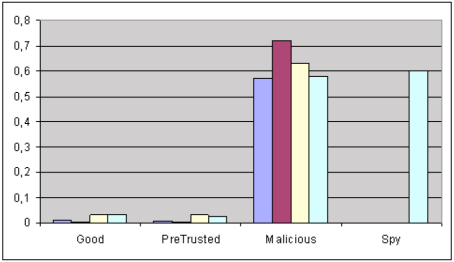
The badness is able to differentiate between good and malicious peers but it does not help in discovering spies.

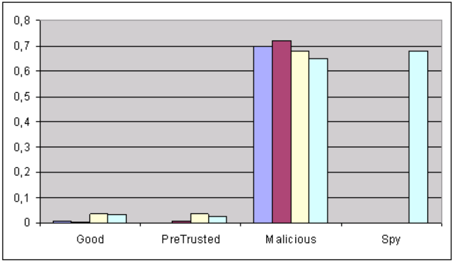
We measure **dishonesty**:

$$\mathbf{dishonesty_i} = \sum_{j \in P} \mathbf{negT_j}$$

where P is the set of peers that i have given positive ratings
The dishonesty is high for all those peers which give good ratings to peers with high badness.

Average dishonesty for models A-D





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Settings

- 100 good peers
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- probability to supply corrupted files equals to 2% for good peers

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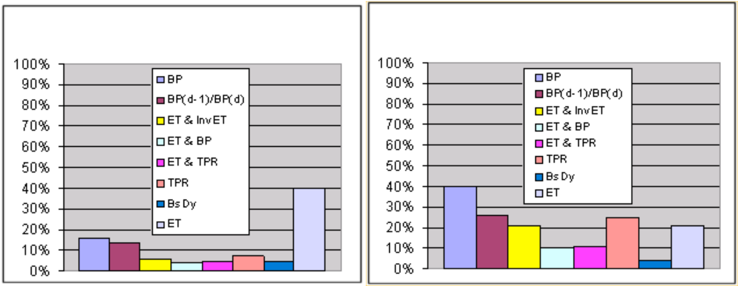
Evaluation

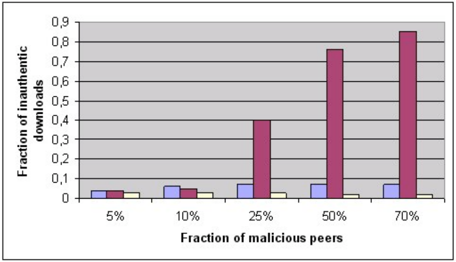
We consider the average ratio between the number of inauthentic downloads and the total number of downloads

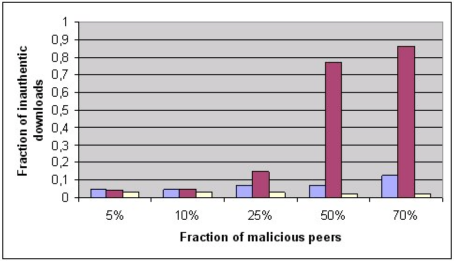
Comparison



Inauthentic downloads for threat model D (malicious and spies) and threat model G (plus smartness)

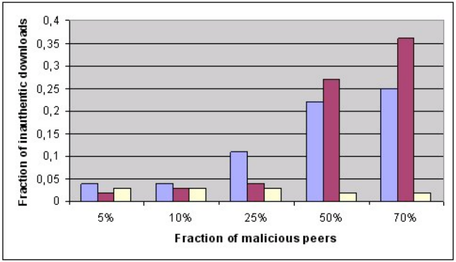
Threat models A (individuals) and B (collective)

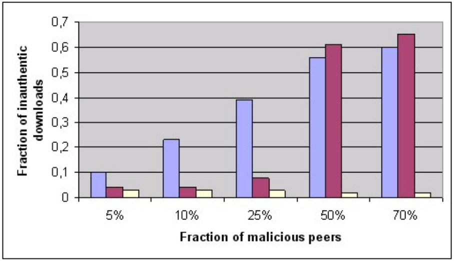




EigenTrust, E. + TruncatedPR, E. + badness + dishonesty

Threat model C (camouflage) and D (spies)





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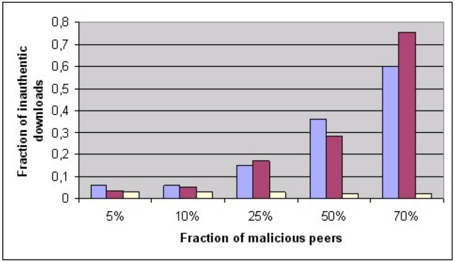
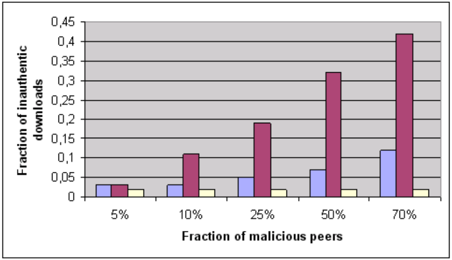
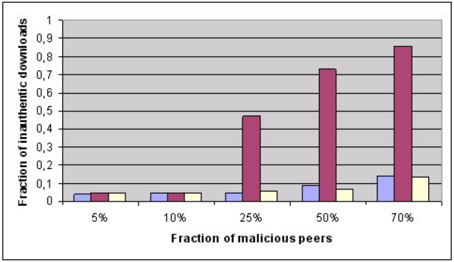
Threat model G (smart) and H (smart+camouflage)

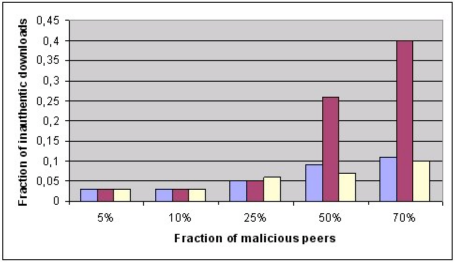




EigenTrust, E. + TruncatedPR, E. + badness + dishonesty

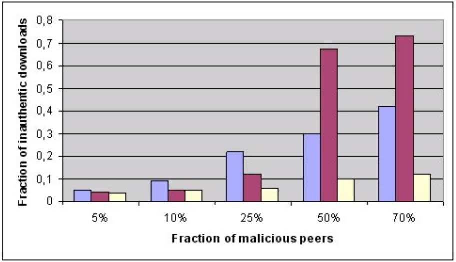
Variant: provide bad files, but be honest

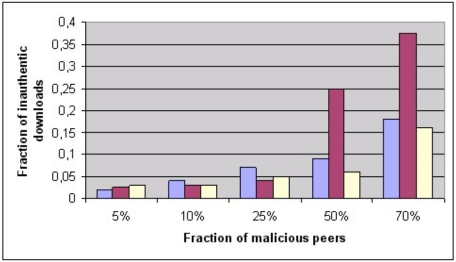




Threat model A',C'

Variant: provide bad files, but be honest; combined attacks





Threat model D+A',D+C'

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Thank you!

http://ewwws.com/pr/przero.php.
PR0 - Google's PageRank 0 Penalty.

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